



## Lecture 5

### Exploiting. Shellcodes (part 2)

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Computer and Network Security  
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push it on the stack and save the pointer

#### Data on stack

```
xor eax, eax  
push eax  
push 0x68732f2f  
push 0x6e69622f  
mov ebx, esp
```

do a jump-call trick (<http://stackoverflow.com/a/15704848>)

#### jump-call trick

```
    jmp MESSAGE      ; 1) lets jump to MESSAGE
GOBACK:
    mov eax, 0x4
    mov ebx, 0x1
    pop ecx           ; 3) we are popping into 'ecx', now we have the
                     ; address of "Hello, World!\r\n"
MESSAGE:
    call GOBACK       ; 2) we are going back, since we used 'call', that means
                     ; the return address, which is in this case the address
                     ; of "Hello, World!\r\n", is pushed into the stack.
    db "Hello, World!", 0dh, 0ah
```

- ▶ stack addresses may differ even if not using ASLR
- ▶ you need a remote connection to send data: netcat, socket API, expect/pexpect API
- ▶ you may need multiple ping-pongs with the remote service
- ▶ pwntools (<https://github.com/Gallopsled/pwntools>) makes it easier

- ▶ strict input validation
- ▶ very limited set of instructions
- ▶ `http://www.phrack.org/issues.html?issue=57&id=15#article`
- ▶ use initial limited shell code to write extended shell code

- ▶ initialize an environment variable with the shellcode string
- ▶ environment variable is placed on the stack of main
- ▶ may be large enough to store large shellcodes
- ▶ unable to be done if stack is non-executable

- ▶ enough to overwrite the code pointer
- ▶ not enough to store the shellcode
- ▶ only use the buffer to overwrite the code pointer
- ▶ place the shellcode in a different location

- ▶ two-phase attack
- ▶ overwrite the code pointer with the address of main (or that of another function)
- ▶ call the vulnerable read/fgets/etc. function again
- ▶ you may use the first call to leak data or make some more room and the second call for the actual attack



- ▶ place the shellcode on the heap
- ▶ requires a heap buffer overflow
- ▶ made difficult by ASLR and non-executable flags

- ▶ uses `printf()` functions that don't do proper checking of arguments
- ▶ may use `%x` and `%s` to read arbitrary data and string from memory
- ▶ may use `%n` to write arbitrary data into memory and possibly trigger a shellcode execution
- ▶ `puts()` may be used; pass an address with information you want to leak

- ▶ if stack is non-executable, one may not execute code on the stack → no shellcode
- ▶ we could call the `system` library call with the `"/bin/bash"` argument
- ▶ with the help of a buffer overflow one overwrites the return address causing a call to `libc`
- ▶ this is restricted to only functions available in `libc`
- ▶ one must know in advance the address of the `system` library call
- ▶ the `"/bin/bash"` may be stored in an environment variable (or is already stored in the `SHELL` environment variable) and its address may be placed on the stack

- ▶ using existing sequences ending in `ret` from the program executable code
- ▶ sequences are programmed on the stack and then executed one by one to provide the required effect
- ▶ sequences are called gadgets
- ▶ we'll talk more about these in the future classes

### Generate shellcode in PEDA

```
gdb-peda$ shellcode generate x86/linux exec
```

- ▶ <https://docs.pwntools.com/en/stable/>,  
<https://github.com/Gallopsled/pwntools>
- ▶ automate exploiting tasks
- ▶ channels
- ▶ ELF inspection
- ▶ return oriented programming
- ▶ shellcodes
- ▶ packing/unpacking

## Skeleton for using pwntools

```
from pwn import *

local = False
if local == True:
    io = process("/path/to/executable")
else:
    HOST = "141.85.100.200"
    PORT = 31337
    io = remote(HOST, PORT)

# TODO: Create shellcode, payload. Do ping-pong with the vulnerable program.
...
```

## pwntools example

```
from pwn import *

io = process("/path/to/executable")

buffer_start = 0x08424242
buffer_to_ret_address_offset = 0x2c

# Craft payload: shellcode + padding + overwrite_address
shellcode = asm(shellcraft.i386.linux.sh())
payload = shellcode + (buffer_to_ret_address_offset - \
    len(shellcode)) * "A" + p32(buffer_start)

# Send payload to overwrite return address with buffer
# start address (buffer stores shellcode).
io.send(payload)

# Do recv if required and other ping-pong with the vulnerable program.
...

# Turn interactive and use the shell.
io.interactive()
```



- ▶ `http://www.metasploit.com/`
- ▶ metasploit framework (open source) + metasploit project
- ▶ penetration testing platform
- ▶ ships with hundreds of exploits (payloads)
- ▶ makes it easy to develop exploits

- ▶ shellcode data
- ▶ jump-call trick
- ▶ alphanumeric shellcode
- ▶ environment variable
- ▶ string format attack
- ▶ return-to-libc
- ▶ pwntools
- ▶ shellcraft
- ▶ data packing
- ▶ pwntools tubes

- ▶ <http://www.blackhatlibrary.net/Category:Shellcode>
- ▶ <http://www.shell-storm.org/shellcode/>
- ▶ <http://www.metasploit.com/>

- ▶ The Ethical Hacker's Handbook, 3rd Edition
  - ▶ Chapter 13 & 14
- ▶ A Guide to Kernel Exploitation
  - ▶ Chapter 1: From User-Land to Kernel-Land Attacks
- ▶ The Art of Exploitation, 2nd Edition
  - ▶ Chapter 0x500. Shellcode
- ▶ Hacking Exposed. Malware and Rootkits
  - ▶ Part II: Rootkits
- ▶ <https://www.win.tue.nl/~aeb/linux/hh/hh-10.html>
- ▶ <https://dhavalkapil.com/blogs/Shellcode-Injection/>
- ▶ Smashing the Stack for Fun and Profit:  
<http://insecure.org/stf/smashstack.html>