



Lecture 4
Exploiting. Shellcodes

Computer and Network Security
October 21, 2019
Computer Science and Engineering Department



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Vulnerabilities and Exploits

Runtime Exploiting

Shellcode

Triggering and Placing Shellcodes

Conclusion

- bugs: misbehaving software
- vulnerability: misbehaviour that can benefit an attacker
- exploiting: turning a vulnerability into an advantage for the attacker
- auditing: analyzing an application to determine its vulnerabilities

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- ▶ poor development process
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- ▶ denial of service
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- ▶ stay connected
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- ▶ alter normal execution pattern
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- preparatory phase
- shellcode
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- ▶ spawn shell using `execve` syscall
- ▶ use `setresuid` to restore root privileges (for `setuid`-enabled programs)
- ▶ port-binding shellcode: create listener socket, accept connections, duplicate file descriptors and spawn shell
- ▶ connect-back shellcode: create client socket and connect to remote listener socket (accessible and controlled by attacker), duplicate file descriptors and spawn shell

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- ▶ may be done in C but it is recommended to do it in assembly
 - ▶ allows shorter shellcodes
 - ▶ complete control over the end result (binary machine code)
- ▶ need to use syscalls for `execve`, `setresuid`, `dup2` and others
- ▶ need to place the `/bin/sh` string in memory (or other strings) and pass it as argument to syscall

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└ Using Syscalls in Linux on x86

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CNSO
Using Syscalls in Linux on x86
- eax stores the syscall number
- ebx, ecx, edx, esi, edi store syscall arguments
- use int 0x80 to issue syscall
- syscall numbers in /usr/include/asm/unistd_32.h

setresuid(0, 0, 0) & exit(1)
1 # Fill eax, ebx, ecx and edx with zeros.
2 xor %eax, %eax
3 xor %ebx, %ebx
4 xor %ecx, %ecx
5 xor %edx, %edx
6 mov $164, %al          # Put 164 (setresuid syscall no) in eax.
7 int $0x80             # Issue syscall: setresuid(0, 0, 0).

1 xor %eax, %eax        # Fill eax with zeros.
2 xor %ebx, %ebx        # Fill ebx with zeros.
3 mov $1, %bl           # Put 1 (EXIT_FAILURE) in ebx (only one
                        # byte).
4 mov $252, %al         # Put 252 (exit_group syscall no) in eax.
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Assembly Wrapper
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Assembly Shellcode Sample
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CNSO Building a Shellcode Wrapper

```

1 ASFLAGS = -march=i386 --32
2 CFLAGS = -Wall -m32
3 LDFLAGS = -m32
4
5 .PHONY: all clean
6
7 all: shellcode-wrapper-exit
8
9 shellcode-wrapper-exit: shellcode-wrapper-exit.o
10
11 shellcode-wrapper-exit.o: shellcode-wrapper-exit.s
12
13 clean:
14     -rm -f shellcode-wrapper-exit shellcode-wrapper-exit.o *~

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Makefile

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CNSO Extracting Hex Data

- ▶ actual shellcode is the machine code instruction
- ▶ use objdump on the object file and process the result
- ▶ use echo -en above to print in binary form

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Using objdump to extract hex data
for i in $(objdump -d module-name.o | tr '\t' ' ' | tr ' ' '\n'
| egrep '^([0-9a-f]2$)'); do echo -n "\x$i"; done
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- ▶ position-independent
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  └ \xb8\x01\x00\x00\x00
> GOOD: xor %eax, %eax + inc %eax
  └ doesn't use null bytes
  └ \x31\xc0\x40
> BAD: mov $100, %eax
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- ▶ shellcode triggering
- ▶ shellcode placing
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